

Workshop in OCL and Textual Modelling

Report on Recent Trends and Panel Discussions

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Abstract This paper reports on the panel session of the 17th Workshop in OCL and Textual Modelling. As in previous years, the panel session featured several lightning talks for presenting recent developments and open questions in the area of OCL and textual modelling. During this session, the OCL community discussed, stimulated through short presentations by OCL experts, proposals for improving OCL to increase the attractiveness of textual modelling.

This paper contains a summary of the workshop from the workshop organisers as well as summaries of two lightning talks provided by their presenters.

Keywords: OCL · textual modelling

1 Introduction

Textual modelling in general and OCL in particular are well established. This year does not only mark the 17th edition of the OCL workshop, it also marks the twentieth anniversary of the first publication of the OCL standard by the OMG [3]. Nevertheless, textual modelling in general and OCL in particular is an active field of research. This year, the Workshop in OCL and Textual Modelling features five regular talks and three lightning talks that covered topics such as the translation of OCL to programming and specification languages, proposals for improving textual modelling languages and their tool support, as well as the development of an OCL benchmark.

The lighting talks at the panel session of the workshop provided a platform for the textual modelling community to discuss and present tools, ideas, and proposals to support textual modelling as well as to shape the future of textual modelling. The following sections, each of them contributed by one expert of the field, discuss the different tools and ideas that were discussed during the panel session.

2 Sometimes Postconditions Do Not Suffice

Martin Gogolla and Antonio Vallecillo

2.1 Non-Determinateness and Randomness in OCL

Recently there have been proposals for incorporating the option to express randomness in OCL [2, 4]. In many modelling and simulation environments, the use of random numbers and probability distributions are used to combine definite knowledge with an uncertain view on the result or the population of a test case. Thus, there is an interest to express such requirements in UML and OCL.

OCL already has operations that possess a flavor of randomness, like the operation `any()`. One could also consider a new collection operation `random()` that randomly chooses an element from the argument collection. Our understanding of such operations is that they cannot be characterized only by ‘traditional’ postconditions. In particular special attention has to be given in order to express the difference between `any()` and `random()`: A ‘traditional’ postcondition would characterize ‘one’ call to the respective operation (for example, with `Set{1..6}->includes(result)`); but these two operations must be characterized by ‘many’ operation calls and a comparison between their actual and their expected results. We show with a small example how such a ‘non-traditional’ postcondition in form of an invariant could look like.

2.2 Formulating Randomness Quality Criteria as an Invariant

Consider the class diagram in Fig. 1 that is intended to model a dice. Every time the operation `random6()` is called it should return a random number between 1 and 6. Our expectation for the operation `any()` would be that it can also return any number between 1 and 6, but that different calls to `any()` always yield the same result. In contrast, different calls to `random6()` should show different results.

The attributes in the class `Dice` (see Lst. 1.1) give a simple measure for the quality of the generated random numbers. Basically the attributes say that the number of tests for `random6()` that have to be performed is `numChecks` and that, for example, the difference between (a) the amount of operation calls yielding 2 and (b) the amount of operation calls yielding 5 is at most `deltaMax`. These requirements are formulated as an OCL formula in terms of an invariant of the class `Dice`. The requirement should not be formulated as a `random6()` postcondition because this would lead to a situation where a recursive call to the operation would occur in the postcondition. Much better criteria for the random distribution could be formulated in OCL as well. The purpose of the

```

class Dice
attributes
  numChecks: Integer
  deltaMax: Integer
operations
  random6(): Integer = Set{1..6}->random()
  post returns_1_6: Set{1..6}->includes(result)
constraints
  inv manyRandom6CallsResultInNearlyEquallyDistributedValues:
    -- call random6() many times
    -- store resulting amounts in Sequence{A1,A2,A3,A4,A5,A6}
    -- check differences between A1..A6
  let amts=Set{1..numChecks}->iterate(i: Integer;
    amts: Sequence(Integer)=Sequence{0,0,0,0,0,0} |
    let r=random6() in
    Sequence{
      if r=1 then amts->at(1)+1 else amts->at(1) endif,
      if r=2 then amts->at(2)+1 else amts->at(2) endif,
      if r=3 then amts->at(3)+1 else amts->at(3) endif,
      if r=4 then amts->at(4)+1 else amts->at(4) endif,
      if r=5 then amts->at(5)+1 else amts->at(5) endif,
      if r=6 then amts->at(6)+1 else amts->at(6) endif}) in
    Sequence{1..5}->iterate(i; diffs: Sequence(Integer)=Sequence{} |
      Sequence{i+1..6}->iterate(j; diffs2: Sequence(Integer)=diffs |
        diffs2->including((amts->at(i)-amts->at(j)).abs()))->
        forAll(d | d<=deltaMax)
    end
end

```

Listing 1.1. Specification of the Dice example.

shown invariant is only to demonstrate that many calls to an operation may be necessary in order to express desired properties.

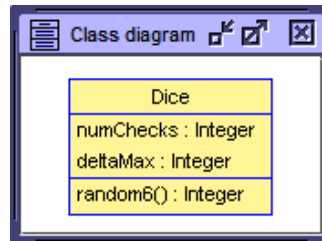


Figure 1. Class diagram for Dice example.

3 Commutative Short Circuit Operators

Edward D. Willink

OCL's 4-level logic has been a source of much unhappiness and while various solutions have been suggested, none have met with enthusiasm. We look at where the unhappiness comes from and thereby suggest a new solution.

The OCL designers defined an underlying model in which all expressions have types. Consequently the mathematical concept of truth was reified by a `Boolean` type with associated Boolean library operations. The designers chose to avoid exceptions. This in combination with UML conformance required a `null` value for the missing value of properties with optional multiplicity, and an `invalid` value for everything bad that might be evaluated.

Unfortunately `null` and `invalid` pollute the simplicity of truths and so the Amsterdam Manifesto [1] elaborates Boolean operators with short-circuit like functionality for problems such as:

```
a <> null and a.doSomething()
```

However the operators remain commutative and so it is suggested that all terms are evaluated in parallel until the result is knowable. A Karnaugh Map defines the mapping from the `true` (T), `false` (F), `null` (ϵ) and `invalid` (\perp) values of `Left` and `Right` inputs to the `and` output.

Left	Right	and	requires	'and2'
T	T	T	T	T
T	F	F	F	F
T	\perp, ϵ	\perp	\perp	\perp
F	-		F	
F	T,F	F		F
F	\perp, ϵ	F		\perp
\perp, ϵ	-		\perp	
\perp, ϵ	T,F, \perp, ϵ	\perp		\perp

Parallel execution is an implementation nightmare and the intermediate `invalid` results can be inefficient. If we eliminate commutative short circuits, we find that `invalid` results are exceptional rather than normal.

```
a <> null requires a.doSomething()
```

A new `requires` operator imposes a left argument first evaluation order for `and`. This avoids the spurious `invalid` results from the right argument and clearly indicates the intent to handle non-truths. The `and` operator can then be used for truths only. Once static analysis verifies that neither left nor right input of an `and` operator can be `null` or `invalid`, an implementation may implement a regular `'and2'` operation that returns `invalid` for any `null` or `invalid` input.

A new `obviates` operator is also needed to regularize `or` short circuiting.

4 Conclusion

The lively discussions both during the lightning talks as well as for each paper that was presented showed again that the OCL community is a very active community. Moreover, it showed that OCL, even though it is a mature language that is widely used, has still areas in which the language can be improved. We

all will look forward to upcoming version of the OCL standard and next year's edition of the OCL workshop.

Acknowledgments. We would like to thank all participants of this years OCL workshop for their active contributions to the discussions at the workshop. These lively discussions are a significant contribution to the success of the OCL workshop series.

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